CIS 313: Intermediate Data Structure

sixth week of the term
red-black trees

1. every node is either red or black
2. the root is black
3. every leaf (null) is black
4. if a node is red, both of its children are black
5. for each node, all simple paths from the node to descendant leaves contain the same number of black nodes
red-black tree height

• (too) simple analysis:
  • the black-height is at most $\log_2 n$
  • the actual height is at most twice the black height
  • so total at most $2\log_2 n$
  • OK, text says at most $2\log_2(n+1)$
red-black tree insertion

• to insert new key $x$
• as always, search to the bottom of the tree for where $x$ would go
• put $x$ there and color it red (to maintain black-height)
• this might cause a problem: two reds in a row
• if no such problem, then done
• if double-red problem, then fix using
  • color shifts or
  • rotation

example: insert 7
RB-INSERT-FIXUP

- section 13.3 of text
- this deals with the double red case after an insertion
- let y be the current node, both it and its parent are red
- let z be the “uncle” of y: the sibling of y’s parent’s parent
- two cases:
  - z is red
    - color shift
    - then check again for double red, possibly continue
  - z is black
    - rotate
    - done
z is red

This swaps a black and red level, preserving black height along these paths, but may create another double-red at the new y.
z is black

rotation (double)

done now
z is black (again, different case)

rotation (single)

there are two other cases similar to these needing single and double left rotations
turn a binary search tree into a red-black tree

```java
if n is root,
    n.color = black
    n.black-quota = height n / 2, rounded up.
else if n.parent is red,
    n.color = black
    n.black-quota = n.parent.black-quota.
else (n.parent is black)
    if n.min-height < n.parent.black-quota, then
        error "shortest path was too short"
    else if n.min-height = n.parent.black-quota then
        n.color = black
    else (n.min-height > n.parent.black-quota)
        n.color = red
    either way,
    n.black-quota = n.parent.black-quota - 1
```
sample BST
example insertions

after insertion of 1, 2, 3, 4, 5, 6 into empty RB tree

let’s continue with 7, 8, …
insert 7

single left rotate
insert 8

- Color shift
- Single left rotate (at 2)
insert 9
insert 10
insert 10 (cont’d)

Note: 4 as root gets colored black at the end
RB Deletion

• BST Deletion Revisited: delete z
  • If z has no children, then just remove it
  • If z has only one child, then splice out z
  • If z has two children, then:
    • Find its successor y
    • Splice out y
    • Replace z’s value with y’s value

-> so the physical node deleted is z in the first two cases and y in the third case
RB Deletion

• Delete z as in BST
• If z has two children, when replace z’s value with the successor’s value, keep z’s color (don’t change z’s color)
• Let y be the node being removed or spliced out in this procedure (y would be either z or successor of z, thus y has at most one child)
• If y is red, no violation of the red-black properties, done
• If y is black, some violations might arise and we need to restore the red-black properties
RB Deletion: y is black

• Let x be the child of y before it was spliced out
  So x is either nil (a leaf) or the only non-nil child of y
Restoring RB Properties

• The RB-DELETE-FIXUP routine in the text, applied to x
• If x is red, so easy, just change its color to black and done
• If x is black:
  • Transform the tree and move x up, until:
    • x points to a red node, or
    • x is the root
  • At each step:
    • need to consider 8 cases; four when x is a left child and four when x is a right child.
    • due to the symmetry, just consider the 4 cases when x is a left child here
  • **REMEMBER:** set the color of x to black in the end
Restoring RB Properties: \( x \) is black and is a left child

x’s sibling \( w \) is red:
Left rotate \( D \), switch colors of \( B \) and \( D \)

x’s sibling \( w \) is black; both \( w \)’s children are black:
Move \( x \) up, change \( w \)’s color to red

x’s sibling \( w \) is black; \( w \)’s right child is black, left child is red:
Left rotate on \( C \), switch colors of \( C \) and \( D \)

x’s sibling \( w \) is black; \( w \)’s right child is red:
Left rotate on \( B \), switch colors of \( B \) and \( D \), change \( E \)’s color to black

The nodes with \( c \) or \( c' \) can be either red or black
example: delete 4

replace 4 with its successor 5, remove 5’s node (y), x is a nil node (child of 5) and it’s black
delete 4 (cont’d)

case 1: left rotate 8, switch colors of 6 and 8
delete 4 (cont’d)

case 2: change color of 7 to red, move up x to 6

x is now red and we are done with the moving up of x

FINAL STEP: change color of x to black
delete 4 (done)

NEXT: delete 1 from this
y is 1 and x is the nil child of 1 (black)
deleting 1 from previous

case 2: change color of 3 to red and move up x to 2
delete 1 (cont’d)

case 2 again: change color of 8 to red, move x up to 5

done since x is at root (we’ve reduced the black height of the tree)
remove 2

y is 2, x is child of y (3), x is red, so just need to change color of x (3) to black

NEXT: delete 3 from this
y is 3, x is the nil node (child of y)
x is black
removing 3 from previous

case 1: left rotate 8, switch colors of 5 and 8
set up as case 4:

motivated by a “transfer” in 2-3-4 tree

case 4: left rotate B (5), switch colors of 5 and 6, change color of 7 to black
Done!
remember: all cases come with mirror image

• here x is right child of parent
• the left child of w is red
• fix-up can be completed with a right rotation and color changes
• note that the blue nodes (B and C) can be either red or black
B-trees

• very important data structure in computer science
• database indexing, hard disk referencing, MongoDB, ...
• balanced, multi-way search tree
• many slight variations, we will use definition in CLRS text
• idea is that nodes are large and fit into a disk block (minimum amount of data that’s pulled off a hard drive)
• node size parameters (here called t) depend on disk speeds, block sizes, etc.
B-tree specifications

- fixed parameter $t$, called *minimum degree*
- nodes have between $t-1$ and $2t-1$ keys
- so therefore they have between $t$ and $2t$ children
- root is exception: it may have as few as 1 key (2 children)
- all null pointers have the same depth (distance from root)
- a 2-3-4 tree is a B-tree with minimum degree $t=2$
different texts: things to look for

- top-down versus bottom-up insertion
  - CLRS does top-down, split full nodes during search
  - unlike how it does RB trees
  - bottom-up more common in practice, less wasted space
- ties to left or right
  - no duplicates here
  - need to know for B+ trees, which have all keys at an additional “leaf level”
- left/right bias: if middle key not well defined (when splitting a node with even number of keys)
B-tree node format

• each node can have between t and 2t children
• a node might look like
  • \(<P_0, K_1, P_1, K_2, P_2, \ldots, P_{q-2}, K_{q-1}, P_{q-1}>\)
  • for t <= q <= 2t (except for root 2<=q<=2t)
• during a search, we split full nodes
• a node is full when it has 2t-1 keys
node split (shown for t=3)

split a full node

becomes

promote c by inserting into parent
B-tree height

• theorem 18.1: if \( n \geq 1 \), then for any B-tree containing \( n \) keys of height \( h \) and minimum degree \( t \geq 2 \), \( h \leq \log_t \frac{n+1}{2} \).

• example: \( t=50 \) and \( n=100,000,000 \)
  • \( h \leq \log_{50} 50,000,000.5 \approx 4.53 \leq 5 \)

• suppose 20 records fit on a page

• without the index to find an item we’d need to search about half the \( 100,000,000/20=5,000,000 \)

• with the index we need at most 5+1=6 disk accesses (5 for the tree nodes and one for the page containing that key’s record)
exercise 18.2-1

insert into initially empty B-tree of min degree t=2 the key values F, S, Q, K, C, L, H, T, V, W, M, R, N, P, A, B, X, Y, D, Z, E

after the first 3 values:

<F, Q, S>

a search to place K causes a split:

<after which K is placed:>

<F, K> <S>
insert into initially empty B-tree of min degree $t=2$ the key values $F, S, Q, K, C, L, H, T, V, W, M, R, N, P, A, B, X, Y, D, Z, E$

place C

search for $L$ splits full node:

split

insert

insert into initially empty B-tree of min degree $t=2$ the key values $F, S, Q, K, C, L, H, T, V, W, M, R, N, P, A, B, X, Y, D, Z, E$
insert into initially empty B-tree of min degree t = 2 the key values F, S, Q, K, C, L, H, T, V, W, M, R, N, P, A, B, X, Y, D, Z, E

place H, T, V

insert W

search

split

insert
insert into initially empty B-tree of min degree $t=2$ the key values
ETC....
aside: splay trees

- maintain balance kind of
- insertion, deletion, union take $O(lg\ n)$ amortized time
  - a series of $m$ of these operations on $n$ keys takes total time worst case $O(m*lg n)$
  - possible that one operation takes $O(n)$ time but cannot happen often
- NO balance information needs to be stored at node (balance factor, color)
- used in DNS servers sometimes
splay trees use rotations

• idea is that whenever a node is accessed, it is moved to the root by a series of rotations
  • LOTS OF ROTATIONS!
  • and slightly different ones
• if x is child of root, it is moved upwards with a single rotation
  • called a ZIG rotation
• if x has a (RL or LR) grandparent, it is moved up with a double rotation
  • called a ZIG-ZAG rotation
• if x has a (LL or RR) rotation, then moved up with a special rotation
  • ZIG-ZIG
• idea: zig-zigs and zig-zags tend towards rebalancing
zig-zig
zig-zag
example: find 7
zig-zig, and done