Main topics of the week:
- Intro to Programming Language Concepts
- Programming Language Domains, Paradigms
- Compilation versus Interpretation

Motivation for Studying Programming Languages
This course will introduce you to the fundamental ideas used in the design of programming languages. It is not a course about how to program in a particular language and we assume you have some experience programming in an imperative language like C, C++, or Java. This is not a course in how compilers work or how to design a compiler. Nor is it a course in the theory of computing. There are separate courses in Compilers and Automata Theory that cover those subjects in detail. We will see some things like grammars and memory management that certainly overlap with those other courses, but the focus here is on the general aspects of programming languages. We will introduce you to some new programming paradigms and languages, but you will not get a complete treatment of these languages. Rather, you will develop an appreciation for the differences between languages and increase your ability to teach yourself about new languages. And, you will understand the tradeoffs between characteristics of programming languages so that you can make better choices about what language is most appropriate for a given application, and what are the important issues in designing a new language.

One of the first questions to ask, is what drives the creation of a programming language in the first place? A computer is just a machine that is capable of executing instructions. A “program” is just the set of instructions. So we can take as a definition:

A **programming language** is a notational system for describing computation in machine readable and human readable form.

At the machine level, these instructions are just code – they constitute a “language” that the machine can understand, e.g., fetch a value from a memory location, add one value to another, store a value. However, if you have looked at any machine microcode or even assembly language, you know it is not easy for most people to understand. After all, it reflects the low level at which the machine operates, and when humans design solutions for problems, these solutions are typically thought of at a “higher” level. So long ago, higher level languages were developed as the preferred way of programming. Machine code is still there, since that is the fundamental way of directing the computer, but programming as a problem solving activity is done at a higher level which is eventually translated mechanically to the machine level. The first level up from the machine microcode is an assembly language. Assembly language maps directly to the machine code, but is more symbolic, rather than just the binary strings that the machine really understands. Programming assembly language is certainly an improvement over programming microcode, but still forces the programmer to deal with the machine’s model of computation. During the 1960’s, Fortran (Formula Translator) was created and became popular. The successful characteristic of Fortran was that it was a programming language that looked a lot like the symbolic algebraic formulas that were being used as the notation in describing the problems people were trying to solve with computers. At the time, there was great doubt about whether this was a good thing to do, since it was considered doubtful that a translator could produce reasonable code. And of course, as it turns out, the translators were able to produce reasonable code, enough so that the
advantage of programming in a language closer to the solution domain far outweighed any inefficiency of the translated code over hand crafted microcode. Fortran became very successful for numerical applications, and Algol attempted to unify various similar languages for these applications. At the same time, Lisp gave rise to a strain of functional programming languages, a fundamentally different paradigm. With the advent of the minicomputer toward the end of the 1960’s, the C programming language, sharing much of the structure of Fortran but with more general purpose computing in mind, was designed along with the Unix operating system and C became one of the most widely used programming languages. ML merged ideas from both these strains. In the 1980’s, object oriented programming came into vogue and C++ and Java are the main successful languages (as measured by adoption) of that type. You can consult the genealogy in the Louden text for a more complete picture of the history of the programming languages.

However, what really drives the development of programming languages? The basic goal is just to perform some type of computation. A programming language is created to:
1) Make programming easier for people
2) Ensure more correct, bug free programs
3) Make efficient use of the computer resource
Sometimes these goals are at odds – a language that ensures that only “perfect” programs can be written may not produce very efficient code. Although one could argue that efficiency is the domain of interpreters and compilers, and should be totally transparent to the programmer, the reality is that the design of a language has a lot to do with whether we can implement compilers and interpreters to produce efficient code. Some folks would even argue that the implementation of the language is the driving force for the language design. Likewise, the design of the language has a lot to do with whether we can guarantee the “correctness” of the program. And of course, the design of a language has everything to do with whether the programming task is more convenient, i.e., easier. Finally, the programming language conveys a certain aesthetic about a program – mostly our concern about this is that the language makes a program easier to read and understand.

Since there is basically one fundamental computing model (as expressed at the most basic level in the Turing Machine), any programming language can theoretically be used to perform any programming task. After all, from an execution standpoint, the original language is irrelevant – a compiler, written in some language, just reads another language and produces some set of machine instructions. The machine doesn’t really “care” where those instructions came from – all of this programming language concern is really for us poor humans.

Toward the goal of making programming easier for people (i.e., human readable), programming languages deal with abstractions that keep the low level computing model at a greater distance. Data abstractions give us a higher level approach to thinking about the data so that we do not get caught up in the machine details. For example, when we write a program dealing with numbers, we don’t really want to get bogged down in how arithmetic is performed – e.g., two’s complement versus one’s complement, etc. The details are clearly important to the actual computation, and the machine designer must be very concerned with understanding this and getting it right, but we would like to stay above this level of detail. That is, we want to deal with the abstraction of a counting number. Moreover, we would like to refer to number constants with our usual
mathematical notation, rather than specifying bit patterns, and we would like to refer to
number entities symbolically – i.e., we want to use variables, which is another data
abstraction. The ability to lay out structures and classes to collect our data together is
another form of data abstraction. Control abstractions give us a higher level approach
to thinking about the actions performed by the machine. Again, the machine designer
must be concerned with things like how one value is copied to another location by the
right combination of fetch and store instructions, but we would like to view this as the
abstraction of assignment. Similarly, the program counter is used to know which
instruction is to be executed next, but we want to deal at the abstract level of sequential
execution with selection (if-else) and loops. Control can be further structured with the
abstractions of procedures and functions or methods.

Programming Domains
The domain of the applications that will be written in the language heavily influences
the design of a programming language. A general-purpose language is by definition not
domain specific, and in principle does not favor any particular type of application.
However, designing a language for specific types of applications will ideally result in a
language that is more expressive, easier to program, less error prone, and more efficient
for the applications for which it is intended. Some applications have special requirements
(e.g., extremely high need for fault tolerance as in air traffic control or life support
equipment) and the design of the language can facilitate these needs. We could roughly
divide application domains into several areas:
- Scientific applications
- Business applications
- Artificial Intelligence
- Systems Programming
- Scripting Languages
Each of these domains has its own characteristics and requirements. For example,
calculation expressiveness and efficiency may be more important, or parallel processing
may be the highest priority. Scripting languages allow programs to be written very
quickly, but with little attention to efficiency.
Other important characteristics to consider in a programming language are to find the
right balance between complexity and simplicity, between ease of coding and correctness
of the code. For example, we want to be able to express complex control flow, but not to
the extent that the language has so many keywords and nuances that it becomes
impossible to predict how a program works.

Evaluating Languages
We have said that the primary goals of a programming language are to make easier,
safer, and more efficient use of the computer. What are the things we should look for to
compare one language to another? One way to evaluate a language (relative to the
programming domain, of course) is to consider the language’s readability, writability,
and reliability. Readability of a language is a measurement of how easy it is to
understand programs written in the language. This is important for maintenance and
enhancement of software. Experience in software engineering indicates that the majority
of the software life cycle is spent on finding bugs and maintenance rather than original
development. So having a language that is readable makes this part of the life cycle easier and less costly. Keeping a language simple tends to make it more readable, but the right balance must be struck between simplicity and expressiveness. Better languages tend to have a number of constructs that can be combined orthogonally. This allows expression of complex ideas without so many different constructs that it is difficult to remember them all. Notation that is too complex or has too many keywords gets in the way of readability. Likewise the organization of the code should be as natural as possible so to promote easy comprehension. Lots of features are nice, but too many, or unexpected hidden behaviors also impair readability. Language design requires a trade off between expressive power and the ability to create unintelligible programs. Designing the language so that the programmer cannot misuse the constructs (e.g., operator overloading) deprives the language of expressive power. The syntax and grammar of a language can facilitate or improve readability.

While readability contributes to easy maintenance of software, writability is a general measure of the ease of original software development. Many of the same factors that contribute to readability also contribute to writability, but in addition, language features that support data abstraction and constructs that give greater expressivity help to make a language writable. Some of the latter may be at odds with readability, e.g., the compact notation possible in expressions in the C language (with side effects, prefix/postfix operators, etc.) can allow a lot to be done with a small amount of code, but may make the code incomprehensible.

Finally, the reliability of a language generally refers to how much analysis can be performed by the compiler or interpreter to ensure that the design of the software is using the language constructs correctly. Type checking is a static analysis that can be performed to ensure types are used correctly, and as we will see later, languages can have various degrees of the strength of type checking, ranging from typeless languages to very strongly typed languages. The language may also include features to promote runtime reliability, such as exception handling that allows cleaner software error recovery, or memory management and garbage collection that avoids misuse of program storage.

Design Considerations

In addition to the general criteria above for evaluating languages, there are some other criteria that factor into language design. In the early days of languages, execution efficiency was the most important consideration. With the spectacular increases in processing speed, execution efficiency is not as critical, but it is still important since we keep designing more complex programs that push the computational limits. Efficiency is a more general concern – one aspect is the execution speed, but there are also efficiency aspects to memory use, programming design speed/ease, maintenance ease, compilation speed, etc. For example, static language constructs permit static analysis, which may not only speed up compilation, but may also speed up execution by avoiding the need for certain runtime checks. Lack of automatic garbage collection may slow down the execution of a program, but may speed up its development. So often times we have conflicts between various efficiency aspects and other language criteria like readability and writability.
Regularity is a measure of how well a language integrates its features, so that there are no unusual restrictions, interactions, or behavior. The constructs should be general enough to accommodate most programming needs, but not so general as to make the language unwieldy. Constructs should typically be orthogonal – there should not be unexpected interactions between different constructs. This makes it easier to learn and use each construct correctly. The language should be uniform in that things that look the similar should behave similarly, and different constructs should really be different and look different. In C, functions are not general since there are no local (nested) functions, yet there are global and local variables. Declarations are not uniform in that class declarations must be followed by a semi-colon, but function declarations should not be. Parameters are not orthogonal to type in Java – primitive variables have copy semantics, but object values have reference semantics (just like arrays versus other types in C/C++).

Simplicity as a design goal means to make things as simple as possible, but not simpler. Ideally, this improves readability, writability, and probably efficiency and reliability, but while C is a simple language, one could argue about how well it meets these other measures. Likewise, expressiveness makes it possible to express conceptual abstractions directly and simply and furthers these design goals. We can also design a language to be extensible and allow the programmer to extend the language in various ways (C++ operator overloading). Security as a language characteristic means that programs cannot do unexpected damage (Java has this as a prime design feature). A good language should have a precise definition that can answer programmers and implementers’ questions. (Most languages today have precise definitions, but only ML has a mathematical definition.) Being able to run the same on any machine is machine-independence and Java is the prime example of a language designed for this purpose (write once, run anywhere). Restrictability is the property that a programmer can program effectively in a subset of the full language. This was Stroustrup’s primary design imperative for C++: you shouldn’t have to pay runtime penalties for features you don’t use.

Programming Paradigms
In this course, we will look at four basic paradigms for programming languages. These are not the only ways of looking at programming languages, but are a convenient way of roughly categorizing languages by their characteristics. Some languages share characteristics from several paradigms. The four paradigms are: imperative, object oriented, functional, and logic.

The main characteristic of an imperative language is that it is closest to the architecture of the (von Neumann) machine, and is execution oriented and allows assignment. An imperative language basically specifies a sequence of state changing actions. Variables are used to refer to the abstract machine’s memory locations, and the contents (state) changes as the sequence of actions unfolds. The key operation in an imperative language is assignment. Examples are Fortran, Algol, and C.

In a functional language, there are no named memory locations – everything is a function call, and values are passed to functions and returned by functions. In particular, there is no assignment. The key operation in a functional language is function application, i.e., calling a function. Recursion figures prominently in the use of a
functional language. Examples are ML and Scheme. The functional paradigm comes from mathematics, and allows programs to be defined more precisely and reasoned about more easily.

Logic languages use a formal logical specification of a problem. These specifications will indicate how to recognize a solution, but not how to find it. The solution will be arrived at through a reasoning process. The key operation is unification. Prolog is a logic programming language. The logic paradigm also has its roots in mathematics – symbolic logic.

In object oriented languages, the focus is on the data abstraction rather than the execution. These data objects communicate with each other and the key operation is this message passing. Object oriented languages may be imperative or functional. Examples are C++, Java.

In addition to considering the different programming paradigms, we also consider the differences between interpreted languages and compiled languages. The compiler statically analyzes program “source” and produces code that is then run on the computer. The compiled program can be executed many times. An interpreted language is dynamically interpreted as it is read and executed by the interpreter. This is for a single execution, and the interpretation of the program source must be done again for all subsequent executions. The latter approach is usually not as efficient, but can more accurately understand the intent of the programmer, as well as allow the program to change dynamically during execution.