CIS 631
Parallel Processing

Lecture 12: Parallel Performance Analysis

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Acknowledgements

☐ Portions of the lectures slides were adopted from:


Outline

- Review
- Parallel performance analysis problem
- Parallel performance analysis methodology
- Measurement and analysis techniques
Parallel Programming

- To use a scalable parallel computer, you must be able to write parallel programs.
- You must understand the programming model and the programming languages, libraries, and systems software used to implement it.
- Unfortunately, parallel programming is not easy.
Parallel Programming Models

- Two general models of parallel program
  - Task parallel
    - Problem is broken down into tasks to be performed
    - Individual tasks are created and communicate to coordinate operations
  - Data parallel
    - Problem is viewed as operations of parallel data
    - Data distributed across processes and computed locally

- Characteristics of scalable parallel programs
  - Data domain decomposition to improve data locality
  - Communication and latency do not grow significantly
Shared Memory Parallel Programming

- Shared memory address space
- (Typically) easier to program
  - Implicit communication via (shared) data
  - Explicit synchronization to access data
- Programming methodology
  - Manual
    - Multi-threading using standard thread libraries
  - Automatic
    - Parallelizing compilers
    - OpenMP parallelism directives
  - Explicit threading (e.g. POSIX threads)
Parallel Programming Model: Threads

- Global style
- Shared and private data
- Work distribution onto threads for global operations
- Domain decomposition determines work distribution
Distributed Memory Parallel Programming

- Distributed memory address space
- (Relatively) harder to program
  - Explicit data distribution
  - Explicit communication via messages
  - Explicit synchronization via messages
- Programming methodology
  - Message passing
    - Plenty of libraries to chose from (MPI dominates)
    - Send-receive, one-sided, active messages
  - Data parallelism
  - Shared virtual memory
Paralle Programming Model: Message Passing

- Local style
- Domain decomposition leads to data distribution
- Explicit communication and synchronization
- Higher programming overhead
- Message passing libraries
Basic Parallel Programming Paradigm: SPMD

- SPMD: Single Program Multiple Data
- One program executes on all processors
- Basic paradigm for implementing parallel programs
- Process-dependent cases are handled inside the program

```plaintext
if (processor == 42) then
    call do_something()
else
    call do_something_else()
endif
```

- Parallelism is “programmed in”
- Easier to manage program for scalability
Parallel Programming: Still a Problem?
Parallel Computing and Scalability

- Scalability in parallel architecture
  - Processor numbers
  - Memory architecture
  - Interconnection network
  - Avoid critical architecture bottlenecks
- Scalability in computational problem
  - Problem size
  - Computational algorithms
    - Computation to memory access ratio
    - Computation to communication ratio
- Parallel programming models and tools
- Performance scalability
**Amdahl’s Law**

- $T_{seq}$: sequential execution time that cannot be parallelized
- $T_{par}$: sequential execution time that can be parallelized
- $T_1 = T_{seq} + T_{par}$  □  $T_{par} = T_1 - T_{seq}$
- $T_p = T_{seq} + T_{par} \div p$ (assume fully parallelized)
- As $p \to \infty$, $T_p \to T_{seq}$
- Let $f_{seq}$ be the fraction $T_{seq} / T_1$ and $S_p = T_1 / T_p$
- Speedup $= S_p = T_1 / T_p = T_1 / (T_{seq} + T_{par} / p)$
  $= 1 / (f_{seq} + T_{par} / pT_1) = 1 / (f_{seq} + (1 - f_{seq}) / p)$
  □ As $p \to \infty$, $S_p = S_\infty \to 1 / f_{seq}$
- Speedup bound is determined by the degree of sequential execution time in the computation, not # processors!!!
Amdahl’s Law and Scaled Speedup

- Amdahl’s Law makes it hard to obtain good speedup

<table>
<thead>
<tr>
<th>$f_{seq}$ * 100%</th>
<th>10%</th>
<th>5%</th>
<th>2%</th>
<th>1%</th>
<th>.1%</th>
</tr>
</thead>
<tbody>
<tr>
<td>$S_\infty$</td>
<td>10</td>
<td>20</td>
<td>50</td>
<td>100</td>
<td>1000</td>
</tr>
</tbody>
</table>

- Change perspective on the problem
  - Consider scaling of problem size as # processors scale

- $T_{seq}$ : sequential execution time (1 and $p$ processors)
  - $T_p = T_{seq} + T_{par}$, $T_1 = T_{seq} + pT_{par}$

- Let $f_{par}$ be the fraction $T_{par} / T_p$

- Scaled speedup $= S_p = 1 + (p-1)T_{par} / T_p = 1 + (p-1)f_{par}$
Parallel Performance

- To use a scalable parallel computer well, you must be able to write high-performance parallel programs.
- To get high-performance parallel programs, you must understand and optimize performance for the combination of programming model, algorithm, language, platform,
- Unfortunately, parallel performance analysis and optimization is not an easy process.

Source: Bernd Mohr
Parallel Performance Evaluation

- Study of performance in parallel systems
  - Models and behaviors
  - Evaluative techniques
- Evaluation methodologies
  - Analytical modeling and statistical modeling
  - Simulation-based modeling
  - Empirical measurement, analysis, and modeling
- Purposes
  - Planning
  - Diagnosis
  - Tuning
Performance Observability (My Guiding Thesis)

- Performance evaluation problems define the requirements for performance analysis methods.

- *Performance observability* is the ability to “accurately” capture, analyze, and present (collectively *observe*) information about computer system/software performance.

- Tools for performance observability must balance the *need* for performance data against the *cost* of obtaining it (environment complexity, performance intrusion):
  - Too little performance data makes analysis difficult.
  - Too much data perturbs the measured system.

- Important to understand performance observability complexity and develop technology to address it.
(Parallel) Performance Analysis Process

- Implementation
- Preparation
- Performance Analysis
- Program Tuning
- Production

Refinement

Measurement

Analysis

Ranking
Parallel Performance Analysis Environment

- **Performance Hypothesis**
  - hypothesis refinement from empirical results
  - general performance results add to performance knowledge base

- **Performance Observation**

- **Stored Performance Knowledge**
- **Experimental Performance Data**

- System/program characteristics plus performance knowledge used for initial hypothesis
- Constraints on observational capabilities, invocation of measurement tools
- Analysis, modeling, and presentation of empirical data
- Parallel system
- Parallel program

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Lecture 12
Performance Analysis and Tuning

- Successful parallel performance tuning process
  - **Characterization**: finding critical performance problems
  - **Diagnosis**: determining performance problem causes
  - **Hypothesis testing**: selection of performance optimization
  - **Hypothesis validation**: analyzing tuning results

- Reasoning and intuition only take you so far

- Need to make empirical observations
  - Performance instrumentation tools
  - Performance measurement tools
  - Performance analysis tools
Factors which determine a program's performance are complex, interrelated, and sometimes hidden.

- **Application related factors**
  - Algorithms dataset sizes
  - Memory usage patterns
  - I/O communication patterns

- **Hardware related factors**
  - Processor architecture
  - Memory hierarchy

- **Software related factors**
  - Operating system
  - Compiler preprocessor

- **Performance Factors**
  - Task Granularity
  - Load Balancing
  - Amdahl's Law
  - I/O network
  - Communication protocols
  - Libraries
Utilization of Computational Resources

- Often resources are under-utilized or used inefficiently
- Identifying these circumstances can give clues to where performance problems exist
- Resources may be “virtual” (i.e., not a physical resource)
  - Thread or process
- Performance analysis tools are essential to optimizing an application's performance
  - Can assist you in understanding what your program is "really doing”
  - May provide suggestions how program performance should be improved
Performance Analysis and Tuning: The Basics

- Most important goal of performance tuning is to reduce a program's wall clock execution time
  - Iterative process to optimize efficiency
  - Efficiency is a relationship of execution time
- So, where does the time go?
- Find your program's hot spots and eliminate the bottlenecks in them
  - Hot spot: an area of code within the program that uses a disproportionately high amount of processor time
  - Bottleneck: an area of code within the program that uses processor resources inefficiently and therefore causes unnecessary delays
- Understand what, where, and how time is being spent
Sequential versus Parallel Performance

- Sequential performance is all about how time is distributed and what resources are used where and when.
- Parallel performance is about sequential performance AND parallel interactions.
  - Sequential performance is the performance within each thread of execution (i.e., its sequential performance).
  - Parallel interactions lead to overheads:
    - synchronization
    - communication
  - Parallel interactions also lead to parallelism inefficiency:
    - load imbalances.
Sequential Performance Tuning

- Sequential performance tuning is a *time-driven* process
- Find the thing that takes the most time and make it take less time (i.e., make it more efficient)
- May lead to program restructuring
  - Changes in data storage and structure
  - Rearrangement of tasks and operations
- May look for opportunities for better resource utilization
  - Cache management is a big one
  - Locality, locality, locality!
  - Virtual memory management may also pay off
- May look for opportunities for better processor usage
Parallel Performance Tuning versus Sequential

- In contrast to sequential performance tuning, parallel performance tuning might be described as conflict-driven or interaction-driven.

- Find the points of parallel interactions and determine the overheads associated with them.

- Overheads can be the cost of performing the interactions:
  - Transfer of data
  - Extra operations to implement coordination

- Overheads also include time spent waiting:
  - Lack of work
  - Waiting for dependency to be satisfied
Interesting Performance Phenomena

- Superlinear speedup
  - Speedup in parallel execution is greater than linear
  - $S_p > p$
  - How can this happen?

- Need to keep in mind the relationship of performance and resource usage

- Computation time (i.e., real work) is not simply a linear distribution to parallel threads of execution

- Resource utilization thresholds can lead to performance inflections
How Is Time Measured?

- How do we determine where the time goes?
- “A person with one clock knows what time it is, a person with two clocks is never sure.” Confucious
- “Define time.” Bill Clinton (attributed)
- Time is only as good (accurate) as the clock we use
- Clocks are not the same and, thus, time is not the same
  - Wallclock time – measured against “real” time
  - CPU (virtual) time – time accumulates (i.e., “ticks”) only when process is executing
  - Clocks have different resolutions and overheads for access
    - affects accuracy
## IBM SP Timers

<table>
<thead>
<tr>
<th>Timer</th>
<th>Usage</th>
<th>Wallclock / CPU Time</th>
<th>Resolution</th>
<th>Languages</th>
<th>Portable?</th>
</tr>
</thead>
<tbody>
<tr>
<td>time</td>
<td>shell / script</td>
<td>both</td>
<td>1/100th second</td>
<td>any</td>
<td>yes</td>
</tr>
<tr>
<td>timex</td>
<td>shell / script</td>
<td>both</td>
<td>1/100th second</td>
<td>any</td>
<td>yes</td>
</tr>
<tr>
<td>gettimeofday</td>
<td>subroutine</td>
<td>wallclock</td>
<td>microsecond</td>
<td>C/C++</td>
<td>yes</td>
</tr>
<tr>
<td>read_real_time</td>
<td>subroutine</td>
<td>wallclock</td>
<td>nanosecond</td>
<td>C/C++</td>
<td>no</td>
</tr>
<tr>
<td>rtc</td>
<td>subroutine</td>
<td>wallclock</td>
<td>microsecond</td>
<td>Fortran</td>
<td>no</td>
</tr>
<tr>
<td>irtc</td>
<td>subroutine</td>
<td>wallclock</td>
<td>nanosecond</td>
<td>Fortran</td>
<td>no</td>
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<td>dtime_</td>
<td>subroutine</td>
<td>CPU</td>
<td>1/100th second</td>
<td>Fortran</td>
<td>no</td>
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<td>etime_</td>
<td>subroutine</td>
<td>CPU</td>
<td>1/100th second</td>
<td>Fortran</td>
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<td>mclock</td>
<td>subroutine</td>
<td>CPU</td>
<td>1/100th second</td>
<td>Fortran</td>
<td>no</td>
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<tr>
<td>timef</td>
<td>subroutine</td>
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<td>millisecond</td>
<td>Fortran</td>
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<td>wallclock</td>
<td>microsecond</td>
<td>C/C++, Fortran</td>
<td>yes</td>
</tr>
<tr>
<td>AIX Trace Facility</td>
<td>shell / script / subroutine</td>
<td>wallclock</td>
<td>microsecond</td>
<td>any</td>
<td>no</td>
</tr>
</tbody>
</table>
Next Class

- Parallel performance tools